

# Introduction to Mathematical Finance

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LSSU Math 500

- 1 Optimization Models
  - A Deterministic Optimization Model
  - Probabilistic Optimization Problems

## Subsection 1

# A Deterministic Optimization Model

# The Problem

- Suppose we have  $m$  dollars to invest among  $n$  projects.
- Investing  $x$  in project  $i$  yields a (present value) return of  $f_i(x)$ ,  $i = 1, \dots, n$ .
- The problem is to determine the integer amounts to invest in each project so as to maximize the sum of the returns.
- Let  $x_i$  denote the amount to be invested in project  $i$ .
- The problem (mathematically) is to:

Choose nonnegative integers  $x_1, \dots, x_n$ ,

$$\text{such that } \sum_{i=1}^n x_i = m,$$

$$\text{to maximize } \sum_{i=1}^n f_i(x_i).$$

# Solution Based on Dynamic Programming

- Let  $V_j(x)$  denote the maximal possible sum of returns when we have a total of  $x$  to invest in projects  $1, \dots, j$ .
- $V_n(m)$  represents the maximal value of the problem.
- We determine  $V_n(m)$ , and the optimal investment amounts, by:
  - Finding first the values of  $V_1(x)$ , for  $x = 1, \dots, m$ ;
  - Finding next the values of  $V_2(x)$ , for  $x = 1, \dots, m$ ;
  - $\vdots$
  - Ending with the values of  $V_n(x)$ , for  $x = 1, \dots, m$ .

## Solution Based on Dynamic Programming (Cont'd)

- The maximal return when  $x$  must be invested in project 1 is  $f_1(x)$ .  
So we have

$$V_1(x) = f_1(x).$$

- Suppose that  $x$  must be invested between projects 1 and 2.
  - Let  $y$  be invested in project 2.
  - Then  $x - y$  is available to invest in project 1.
- The best return from investing  $x - y$  in project 1 is  $V_1(x - y)$ .
- So the maximal sum of returns possible when the amount  $y$  is invested in project 2 is

$$f_2(y) + V_1(x - y).$$

- The maximal sum of returns is obtained by maximizing over  $y$ ,

$$V_2(x) = \max_{0 \leq y \leq x} \{f_2(y) + V_1(x - y)\}.$$

## Solution Based on Dynamic Programming (Cont'd)

- In general, suppose that  $x$  must be invested among projects  $1, \dots, j$ .
  - Suppose we invest  $y$  in project  $j$ .
  - Then a total of  $x - y$  is available to invest in projects  $1, \dots, j - 1$ .
- The best return from investing  $x - y$  in projects  $1, \dots, j - 1$  is  $V_{j-1}(x - y)$ .
- So the maximal sum of returns possible when the amount  $y$  is invested in project  $j$  is  $f_j(y) + V_{j-1}(x - y)$ .
- The maximal sum of returns possible is obtained by maximizing the preceding over  $y$ ,

$$V_j(x) = \max_{0 \leq y \leq x} \{f_j(y) + V_{j-1}(x - y)\}.$$

- Let  $y_j(x)$  denote the value (or a value if there is more than one) of  $y$  that maximizes the right side of the preceding equation.
- Then  $y_j(x)$  is the optimal amount to invest in project  $j$  when we have  $x$  to invest among projects  $1, \dots, j$ .

# Solution Based on Dynamic Programming (Cont'd)

- The value of  $V_n(m)$  can now be obtained by first determining  $V_1(x)$ , then  $V_2(x)$ ,  $V_3(x)$ ,  $\dots$ ,  $V_{n-1}(x)$  and finally  $V_n(m)$ .
- The optimal amounts to invest are:
  - $y_n(m)$  in project  $n$ ;
  - $y_{n-1}(m - y_n(m))$  in project  $n - 1$ ;
  - $\vdots$
- This solution approach is called **dynamic programming**.  
It views the problem as involving  $n$  sequential decisions.  
It then analyzes it by determining:
  - The optimal last decision;
  - The optimal next to last decision;
  - $\vdots$



# Example

- Suppose that three investment projects with the following return functions are available:

$$f_1(x) = \frac{10x}{1+x}, \quad x = 0, 1, \dots,$$

$$f_2(x) = \sqrt{x}, \quad x = 0, 1, \dots,$$

$$f_3(x) = 10(1 - e^{-x}), \quad x = 0, 1, \dots$$

- We want to maximize our return when we have 5 to invest.
- We have

$$V_1(x) = f_1(x) = \frac{10x}{1+x}.$$

- Moreover,

$$y_1(x) = x.$$

## Example (Cont'd)

- Now

$$\begin{aligned} V_2(x) &= \max_{0 \leq y \leq x} \{f_2(y) + V_1(x - y)\} \\ &= \max_{0 \leq y \leq x} \left\{ \sqrt{y} + \frac{10(x-y)}{1+x-y} \right\}. \end{aligned}$$

- So we have

$$V_2(1) = \max \left\{ \frac{10}{2}, 1 \right\} = 5, \quad y_2(1) = 0;$$

$$V_2(2) = \max \left\{ \frac{20}{3}, 1 + 5, \sqrt{2} \right\} = \frac{20}{3}, \quad y_2(2) = 0;$$

$$V_2(3) = \max \left\{ \frac{30}{4}, 1 + \frac{20}{3}, \sqrt{2} + 5, \sqrt{3} \right\} = \frac{23}{3}, \quad y_2(3) = 1;$$

$$V_2(4) = \max \left\{ \frac{40}{5}, 1 + \frac{30}{4}, \sqrt{2} + \frac{20}{3}, \sqrt{3} + 5, \sqrt{4} \right\} = 8.5, \quad y_2(4) = 1;$$

$$V_2(5) = \max \left\{ \frac{50}{6}, 1 + 8, \sqrt{2} + 7.5, \sqrt{3} + \frac{20}{3}, \sqrt{4} + 5, \sqrt{5} \right\} = 9, \\ y_2(5) = 1.$$

## Example (Cont'd)

- Continuing, we get

$$\begin{aligned} V_3(x) &= \max_{0 \leq y \leq x} \{f_3(y) + V_2(x - y)\} \\ &= \max_{0 \leq y \leq x} \{10(1 - e^{-y}) + V_2(x - y)\}. \end{aligned}$$

- We compute the values:

- $1 - e^{-1} = 0.632$ ;
- $1 - e^{-2} = 0.865$ ;
- $1 - e^{-3} = 0.950$ ;
- $1 - e^{-4} = 0.982$ ;
- $1 - e^{-5} = 0.993$ .

- So we obtain

$$\begin{aligned} V_3(5) &= \max \left\{ 9, 6.32 + 8.5, 8.65 + \frac{23}{3}, \right. \\ &\quad \left. 9.50 + \frac{20}{3}, 9.82 + 5, 9.93 \right\} = 16.32, \\ y_3(5) &= 2. \end{aligned}$$

## Example (Cont'd)

- Thus, the maximal sum of returns from investing 5 is 16.32;
- The optimal amount to invest in project 3 is  $y_3(5) = 2$ ;
- The optimal amount to invest in project 2 is

$$y_2(5 - 2) = y_2(3) = 1;$$

- The optimal amount to invest in project 1 is

$$y_1(5 - 2 - 1) = y_1(2) = 2.$$

# Concave Return Functions

- A function  $g(i)$ ,  $i = 0, 1, \dots$ , is said to be **concave** if

$$g(i + 1) - g(i) \text{ is nonincreasing in } i.$$

- We will consider concave return functions  $f_i(x)$ .
- This means that the additional (or marginal) gain from each additional unit invested becomes smaller as more has already been invested.

# Solution for Concave Return Functions

- Assume that the functions  $f_i(x)$ ,  $i = 1, \dots, n$ , are all concave.
- Consider the problem of choosing nonnegative integers  $x_1, \dots, x_n$ , whose sum is  $m$ , to maximize  $\sum_{i=1}^n f_i(x_i)$ .
- Suppose that  $x_1^o, \dots, x_n^o$  is an optimal vector for this problem.
- I.e., a vector of nonnegative integers that sum to  $m$ , with

$$\sum_{i=1}^n f_i(x_i^o) = \max \sum_{i=1}^n f_i(x_i),$$

the maximum over all nonnegative integers  $x_1, \dots, x_n$  that sum to  $m$ .

- Now suppose that we have a total of  $m + 1$  to invest.
- We argue that there is an optimal vector  $y_1^o, \dots, y_n^o$  with  $\sum_{i=1}^n y_i^o = m + 1$  that satisfies  $y_i^o \geq x_i^o$ ,  $i = 1, \dots, n$ .

# Solution for Concave Return Functions (Cont'd)

- Suppose we have  $m + 1$  to invest.
- Consider any investment strategy  $y_1, \dots, y_n$ , such that:
  - $\sum_{i=1}^n y_i = m + 1$ ;
  - For some value of  $k$ ,  $y_k < x_k^o$ .
- We have  $m + 1 = \sum_i y_i > \sum_i x_i^o = m$ .
- Hence, there must be a  $j$  such that  $x_j^o < y_j$ .
- Consider the investment strategy that invests:
  - $y_k + 1$  in project  $k$ ;
  - $y_j - 1$  in project  $j$ ;
  - $y_i$  in project  $i$  for  $i \neq k, j$ .
- We argue that this strategy is at least as good as the strategy that invests  $y_i$  in project  $i$  for each  $i$ .

# Solution for Concave Return Functions (Cont'd)

- We must show that  $f_k(y_k + 1) + f_j(y_j - 1) \geq f_k(y_k) + f_j(y_j)$ .
- Equivalently,

$$f_k(y_k + 1) - f_k(y_k) \geq f_j(y_j) - f_j(y_j - 1).$$

- Now  $x_1^o, \dots, x_n^o$  is optimal when there is  $m$  to invest.
- So

$$f_k(x_k^o) + f_j(x_j^o) \geq f_k(x_k^o - 1) + f_j(x_j^o + 1).$$

- Equivalently, we have

$$f_k(x_k^o) - f_k(x_k^o - 1) \geq f_j(x_j^o + 1) - f_j(x_j^o).$$

- Consequently,

$$\begin{aligned} & f_k(y_k + 1) - f_k(y_k) \\ & \geq f_k(x_k^o) - f_k(x_k^o - 1) \quad (\text{by concavity, since } y_k + 1 \leq x_k^o) \\ & \geq f_j(x_j^o + 1) - f_j(x_j^o) \quad (\text{by the preceding inequality}) \\ & \geq f_j(y_j) - f_j(y_j - 1) \quad (\text{by concavity, since } x_j^o + 1 \leq y_j). \end{aligned}$$



# Solution for Concave Return Functions (Cont'd)

- Thus, any strategy for investing  $m + 1$  that calls for investing less than  $x_k^o$  in some project  $k$  can be at least matched by one whose investment in project  $k$  is increased by 1 with a corresponding decrease in some project  $j$  whose investment was greater than  $x_j^o$ .
- Repeating this argument shows that, for any strategy of investing  $m + 1$ , we can find another strategy that:
  - Invests at least  $x_i^o$  in project  $i$ , for all  $i = 1, \dots, n$ ;
  - Yields a return that is at least as large as the original strategy.
- This implies that we can find an optimal strategy  $y_1^o, \dots, y_n^o$  for investing  $m + 1$  that satisfies the inequality claimed.

# Solution for Concave Return Functions (Cont'd)

- We argued that the optimal strategy for investing  $m + 1$  invests at least as much in each project as does the optimal strategy for investing  $m$ .
- It follows that the optimal strategy for  $m + 1$  can be found by using the optimal strategy for  $m$  and then investing the extra dollar in that project whose marginal increase is largest.
- Therefore, we can find the optimal investment (when we have  $m$ ) by:
  - First solving the optimal investment problem when we have 1 to invest;
  - Then solving the optimal investment problem when we have 2 to invest;
  - Then solving the optimal investment problem when we have 3 to invest;
  - $\vdots$

# Example Revisited

- We reconsider the preceding example.
- We have 5 to invest among three projects, with return functions

$$f_1(x) = \frac{10x}{1+x}, \quad f_2(x) = \sqrt{x}, \quad f_3(x) = 10(1 - e^{-x}).$$

- Let  $x_i(j)$  denote the optimal amount to invest in project  $i$  when we have a total of  $j$  to invest.
- We have

$$\max \{f_1(1), f_2(1), f_3(1)\} = \max \{5, 1, 6.32\} = 6.32.$$

- So

$$x_1(1) = 0, \quad x_2(1) = 0, \quad x_3(1) = 1.$$

## Example (Cont'd)

- Now

$$\max_i \{f_i(x_i(1) + 1) - f_i(x_i(1))\} = \max \{5, 1, 8.65 - 6.32\} = 5.$$

- So we have

$$x_1(2) = 1, \quad x_2(2) = 0, \quad x_3(2) = 1.$$

- Further,

$$\begin{aligned} \max_i \{f_i(x_i(2) + 1) - f_i(x_i(2))\} &= \max \left\{ \frac{20}{3} - 5, 1, 8.65 - 6.32 \right\} \\ &= 2.33. \end{aligned}$$

- So we get

$$x_1(3) = 1, \quad x_2(3) = 0, \quad x_3(3) = 2.$$

## Example (Cont'd)

- Continuing,

$$\begin{aligned} \max_i \{f_i(x_i(3) + 1) - f_i(x_i(3))\} &= \max \left\{ \frac{20}{3} - 5, 1, 9.50 - 8.65 \right\} \\ &= 1.67. \end{aligned}$$

- Therefore,

$$x_1(4) = 2, \quad x_2(4) = 0, \quad x_3(4) = 2.$$

- Finally,

$$\begin{aligned} \max_i \{f_i(x_i(4) + 1) - f_i(x_i(4))\} &= \max \left\{ \frac{30}{4} - \frac{20}{3}, 1, 9.50 - 8.65 \right\} \\ &= 1. \end{aligned}$$

- This gives

$$x_1(5) = 2, \quad x_2(5) = 1, \quad x_3(5) = 2.$$

- Thus, the maximal return is

$$6.32 + 5 + 2.33 + 1.67 + 1 = 16.32.$$

# Algorithm

- The following algorithm can be used to solve the problem when  $m$  is to be invested among  $n$  projects, each with a concave return function.
- The quantity  $k$  will represent the current amount to be invested.
- $x_i$  will represent the optimal amount to invest in project  $i$  when a total of  $k$  is to be invested.
  - (1) Set  $k = 0$  and  $x_i = 0$ ,  $i = 1, \dots, n$ .
  - (2)  $m_i = f_i(x_i + 1) - f_i(x_i)$ ,  $i = 1, \dots, n$ .
  - (3)  $k = k + 1$ .
  - (4) Let  $J$  be such that  $m_J = \max_i m_i$ .
  - (5) If  $J = j$ , then  $x_j \rightarrow x_j + 1$ ,  $m_j \rightarrow f_j(x_j + 1) - f_j(x_j)$ .
  - (6) If  $k < m$ , go to step (3).
- Step (5) means that if the value of  $J$  is  $j$ , then:
  - (a) The value of  $x_j$  should be increased by 1;
  - (b) The value of  $m_j$  should be reset to equal the difference of  $f_j$  evaluated at 1 plus the new value of  $x_j$  and  $f_j$  evaluated at the new value of  $x_j$ .

# Remark

- When  $g(x)$  is defined for all  $x$  in an interval, then  $g$  is concave if  $g'(t)$  is a decreasing function of  $t$  (that is, if  $g''(t) \leq 0$ ).
- Hence, for  $g$  concave

$$\int_i^{i+1} g'(s) ds \leq \int_{i-1}^i g'(s) ds.$$

- So

$$g(i+1) - g(i) \leq g(i) - g(i-1).$$

- This is the definition of concavity we used for  $g$  defined on the integers.

# The Knapsack Problem

- Assume we can invest at most  $m$  in the  $n$  projects.
- Suppose one invests in project  $i$  by buying an integral number of shares in that project, with each share:
  - Costing  $c_i$ ;
  - Returning  $v_i$ .
- Let  $x_i$  denote the number of shares of project  $i$  that are purchased.
- Then the problem is to:

Choose nonnegative integers  $x_1, \dots, x_n$ ,

$$\text{such that } \sum_{i=1}^n x_i c_i \leq m,$$

$$\text{to maximize } \sum_{i=1}^n v_i x_i.$$

- We will use a dynamic programming approach to solve this problem.



# The Knapsack Problem (Cont'd)

- Let  $V(x)$  be the maximal return possible when we have  $x$  to invest.
- If we start by buying one share of project  $i$ , then a return  $v_i$  will be received and we will be left with a capital of  $x - c_i$ .
- $V(x - c_i)$  is the maximal return from investing  $x - c_i$ .
- So the maximal return possible if we have  $x$  and begin investing by buying one share of project  $i$  is

$$\begin{aligned} &\text{maximal return if start by purchasing one share of } i \\ &= v_i + V(x - c_i). \end{aligned}$$

- Hence, the maximal return  $V(x)$  that can be obtained from the investment capital  $x$ , satisfies

$$V(x) = \max_{i: c_i \leq x} \{v_i + V(x - c_i)\}.$$

# The Knapsack Problem (Cont'd)

- Let  $i(x)$  denote the value of  $i$  that maximizes  $v_i + V(x - c_i)$ .
- Starting with  $x$ , it is optimal to purchase one share of project  $i(x)$ .
- Starting with

$$V(1) = \max_{i:c_i \leq 1} v_i,$$

it is easy to determine the values of  $V(1)$  and  $i(1)$ .

- This will then enable us to use

$$V(x) = \max_{i:c_i \leq x} \{v_i + V(x - c_i)\}$$

to determine  $V(2)$  and  $i(2)$ .

- And so on.

# The Name

- We introduced the problem:

Choose nonnegative integers  $x_1, \dots, x_n$ ,

$$\text{such that } \sum_{i=1}^n x_i c_i \leq m,$$

$$\text{to maximize } \sum_{i=1}^n v_i x_i.$$

- This problem is called a **knapsack** problem.

It is mathematically equivalent to determining the set of items to be put in a knapsack that can carry a total weight of at most  $m$  when there are  $n$  different types of items, with each type  $i$  item having:

- Weight  $c_i$ ;
- Value  $v_i$ .

# Example

- Suppose you have 25 to invest among three projects whose cost and return values are as on the right.

Project	Cost/Share	Return/Share
1	5	7
2	9	12
3	15	22

$$V(x) = 0, \quad x \leq 4;$$

$$V(x) = 7, \quad i(x) = 1, \quad x = 5, 6, 7, 8;$$

$$V(9) = \max \{7 + V(4), 12 + V(0)\} = 12, \quad i(9) = 2;$$

$$V(x) = \max \{7 + V(x - 5), 12 + V(x - 9)\} = 14, \quad i(x) = 1, \\ x = 10, 11, 12, 13;$$

$$V(14) = \max \{7 + V(9), 12 + V(5)\} = 19, \quad i(x) = 1 \text{ or } 2;$$

$$V(15) = \max \{7 + V(10), 12 + V(6), 22 + V(0)\} = 22, \quad i(15) = 3;$$

$$V(16) = \max \{7 + V(11), 12 + V(7), 22 + V(1)\} = 22, \quad i(16) = 3;$$

$$V(17) = \max \{7 + V(12), 12 + V(8), 22 + V(2)\} = 22, \quad i(17) = 3;$$

$$V(18) = \max \{7 + V(13), 12 + V(9), 22 + V(3)\} = 24, \quad i(18) = 2;$$

and so on.

## Subsection 2

# Probabilistic Optimization Problems

# A Gambling Model with Unknown Win Probabilities

- Suppose that an investment's win probability can be one of three possible values:  $p_1 = 0.45$ ,  $p_2 = 0.55$  or  $p_3 = 0.65$ .
- Suppose also that it will be:
  - $p_1$  with probability  $\frac{1}{4}$ ;
  - $p_2$  with probability  $\frac{1}{2}$ ;
  - $p_3$  with probability  $\frac{1}{4}$ .
- An investor, without any information about which  $p_i$  has been chosen, will take the win probability to be

$$p = \frac{1}{4}p_1 + \frac{1}{2}p_2 + \frac{1}{4}p_3 = 0.55.$$

# Gambling with Unknown Win Probabilities (Cont'd)

- Assume the investor has:
  - Initial fortune  $x$ ;
  - A log utility function.
- By a previous example, we know that the investor:
  - Will invest  $100(2p - 1) = 10\%$  of her fortune;
  - Will have expected utility of her final fortune

$$\begin{aligned}\log(x) + 0.55 \log(1.1) + 0.45 \log(0.9) \\ = \log(x) + 0.0050 = \log(e^{0.0050}x).\end{aligned}$$

# Gambling with Unknown Win Probabilities (Cont'd)

- Suppose now that the investor is able to learn, before making her investment, which  $p_i$  is the win probability.
  - If 0.45 is the win probability, then the investor will not invest. The conditional expected utility of her final fortune will be  $\log(x)$ .
  - If 0.55 is the win probability, the investor will do as shown previously. The conditional expected utility of her final fortune will be  $\log(x) + 0.0050$ .
  - If 0.65 is the win probability, the investor will invest 30% of her fortune. The conditional expected utility of her final fortune will be

$$\log(x) + 0.65 \log(1.3) + 0.35 \log(0.7) = \log(x) + .0456.$$

- Therefore, the expected final utility of an investor who will learn which  $p_i$  is the win probability before making her investment is

$$\begin{aligned} \frac{1}{4} \log(x) + \frac{1}{2}(\log(x) + 0.0050) + \frac{1}{4}(\log(x) + 0.0456) \\ = \log(x) + 0.0139 = \log(e^{0.0139}x). \end{aligned}$$



# An Investment Allocation Model

- An investor has the amount  $D$  available to invest.
- During each of  $N$  time instants, an opportunity to invest will (independently) present itself with probability  $p$ .
- If the opportunity occurs, the investor must decide how much of her remaining wealth to invest.
- If  $y$  is invested in an opportunity then  $R(y)$ , a specified function of  $y$ , is earned at the end of the problem.
- Both the amount invested and the return from that investment become unavailable for future investment.
- The investor's final wealth is equal to the sum of all the investment returns and the amount that was never invested.
- We determine how much to invest at each opportunity so as to maximize the expected value of the investor's final wealth.

# Notation

- Let  $W_n(x)$  denote the maximal expected final wealth when:
  - The investor has  $x$  to invest;
  - There are  $n$  time instants in the problem.
- Let  $V_n(x)$  denote the maximal expected final wealth when:
  - The investor has  $x$  to invest;
  - There are  $n$  time instants in the problem;
  - An opportunity is at hand.

# Determining $V_n(x)$

- Suppose  $y$  is initially invested;
- Then the investor's maximal expected final wealth will be  $R(y)$  plus the maximal expected amount that she can obtain in  $n - 1$  time instants when her investment capital is  $x - y$ .
- The latter quantity is  $W_{n-1}(x - y)$ .
- So the maximal expected final wealth when  $y$  is invested is

$$R(y) + W_{n-1}(x - y).$$

- The investor can now choose  $y$  to maximize this sum,

$$V_n(x) = \max_{0 \leq y \leq x} \{R(y) + W_{n-1}(x - y)\}.$$

# Determining $W_n(x)$

- Suppose the investor has  $x$  to invest.
- Suppose there are  $n$  time instants to go.
- One of the following two cases arises:
  - An opportunity occurs and the maximal expected final wealth is  $V_n(x)$ ;
  - An opportunity does not occur and the maximal expected final wealth is  $W_{n-1}(x)$ .
- Each opportunity occurs with probability  $p$ .
- So we have

$$W_n(x) = pV_n(x) + (1 - p)W_{n-1}(x).$$

# Solution Method

- Start with  $W_0(x) = x$ .
  - We first use the former equation to obtain  $V_1(x)$ , for all  $0 \leq x \leq D$ ;
  - Then use the latter equation to obtain  $W_1(x)$ , for all  $0 \leq x \leq D$ ;
  - Then use the former equation to obtain  $V_2(x)$  for all  $0 \leq x \leq D$ ;
  - Then use the latter equation to obtain  $W_2(x)$ ;
  - $\vdots$
- Let  $y_n(x)$  be the value of  $y$  that maximizes the right side of the former equation.
- The optimal policy is to invest the amount  $y_n(x)$  if:
  - Our current investment capital is  $x$ ;
  - There are  $n$  time instants remaining;
  - An opportunity is present.

# Example

- We work under the following hypotheses:
  - We have 10 to invest;
  - There are two time instants;
  - An opportunity presents itself each instant with probability  $p = 0.7$ , and  $R(y) = y + 10\sqrt{y}$ .
- We find the maximal expected final wealth and the optimal policy.
- We start with  $W_0(x) = x$ .
- We then get

$$\begin{aligned}V_1(x) &= \max_{0 \leq y \leq x} \{y + 10\sqrt{y} + x - y\} \\ &= x + \max_{0 \leq y \leq x} \{10\sqrt{y}\} \\ &= x + 10\sqrt{x}.\end{aligned}$$

- Moreover,  $y_1(x) = x$ .
- Thus,

$$W_1(x) = 0.7(x + 10\sqrt{x}) + 0.3x = x + 7\sqrt{x}.$$

## Example (Cont'd)

- Now we have

$$\begin{aligned}V_2(x) &= \max_{0 \leq y \leq x} \{y + 10\sqrt{y} + x - y + 7\sqrt{x-y}\} \\ &= x + \max_{0 \leq y \leq x} \{10\sqrt{y} + 7\sqrt{x-y}\} \\ &= x + \sqrt{149x},\end{aligned}$$

where calculus gives the final equation, as well as

$$y_2(x) = \frac{100}{149}x.$$

- The preceding now yields

$$\begin{aligned}W_2(x) &= 0.7(x + \sqrt{149x}) + 0.3(x + 7\sqrt{x}) \\ &= x + 0.7\sqrt{149x} + 2.1\sqrt{x}.\end{aligned}$$

## Example (Conclusion)

- Starting with 10, the maximal expected final wealth is

$$W_2(10) = 10 + 0.7\sqrt{1490} + 2.1\sqrt{10} = 43.66.$$

- The optimal policy is to invest:
  - $\frac{1000}{149} = 6.71$ , if an opportunity presents itself at the initial time instant;
  - Whatever of your fortune remains, if an opportunity presents itself at the final time instant.



# Properties of $W_n(x)$ and $V_n(x)$

## Theorem

If  $R(y)$  is a nondecreasing concave function, then:

- (a)  $V_n(x)$  and  $W_n(x)$  are both nondecreasing concave functions;
- (b)  $y_n(x)$  is a nondecreasing function of  $x$ ;
- (c)  $x - y_n(x)$  is a nondecreasing function of  $x$ ;
- (d)  $y_n(x)$  is a nonincreasing function of  $n$ .

- Part (b) states that the more you have the more you should invest.
- Part (c) states that the more you have the more you should conserve.
- Part (d) says that the more time you have the less you should invest each time.