

Read each problem **very carefully** before starting to solve it. Each problem is worth 5 points. It is necessary to show **all** your work. Correct answers without explanations are worth 0 points. GOOD LUCK!!

1. Give the formal inductive definition of the set \mathcal{R} of **regular expressions**:

Basis: $\emptyset, \Lambda, a \in \mathcal{R}$, where $a \in A$;

Induction: If $R, S \in \mathcal{R}$, then

- (a) $(R) \in \mathcal{R}$;
- (b) $R + S \in \mathcal{R}$;
- (c) $R \cdot S \in \mathcal{R}$; and
- (d) $R^* \in \mathcal{R}$.

2. Apply the operator L that associates to a given regular expression the corresponding regular language recursively (**showing all steps**) to discover the regular language $L(b(a^*bc^* + ac))$:

$$\begin{aligned}
 L(b(a^*bc^* + ac)) &= L(b)L(a^*bc^* + ac) \\
 &= L(b)(L(a^*bc^*) \cup L(ac)) \\
 &= L(b)(L(a)^*L(b)L(c)^* \cup L(a)L(c)) \\
 &= \{b\}(\{a\}^*\{b\}\{c\}^* \cup \{a\}\{c\}) \\
 &= \{b\}(\{a^n : n \geq 0\}\{b\}\{c^m : m \geq 0\} \cup \{ac\}) \\
 &= \{b\}(\{a^nbcm : n, m \geq 0\} \cup \{ac\}) \\
 &= \{bac\} \cup \{ba^nbcm : n, m \geq 0\}.
 \end{aligned}$$

3. Write a regular expression for the language L over the alphabet $A = \{a, b, c\}$ consisting of all strings that contain the substring aba and end in c .

The required expression is

$$(a + b + c)^*aba(a + b + c)^*c.$$

4. The equation $(RR)^* = R^*R^*$ between regular expressions is false

Proof:

For alphabet $A = \{a\}$ and $R = a$, the given equation becomes

$$(aa)^* = a^*a^*.$$

Since $a \in L(a^*a^*)$, but $a \notin L((aa)^*)$, we have that

$$L((aa)^*) \neq L(a^*a^*).$$

Therefore, we conclude that

$$(aa)^* \neq a^*a^*.$$